Mind the Gap(s):
Creating & Closing Gaps Between Design and Verification

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Things we will and will not be discussing...

**IS AND IS NOTS...**
This Workshop Will Not...

Remind you of increasing complexities...

...or shorter schedules...

...or fewer resources...

....or closing coverage faster*

...or finding bugs faster*

*at least as the ends, and not the means
This Workshop Will...

Help you prove faster that your design works as intended via mindful attention to gaps in development
Gaps Are Neither Good nor Bad

You succeed or fail based on how you address gaps!

1. Recognize where they occur
2. Align your process to minimize the loss across the gaps
3. Cover the remaining gaps with precision

This workshop will highlight typical gaps, then show examples of:

- **Creating** intentional gaps in development to align work, organizations, skills
- **Closing** gaps precisely
  - That are organizational, geographical, or time-based
  - Between design and implementation
WHERE ARE THE GAPS?

Every point of hand-off is a potential gap in the development process.
Certain Types of Gaps Cause Program Issues

**Documentation**: incomplete, subject to translation, assumptions, interpretation

**Models**: abstractions, inaccurate

**Change**: impacts requirements, constraints, markets, technologies, standards and teams

**Organization**: (and discipline) boundaries create opportunities for information loss

**Knowledge**: need may not align with ability
Asking an OO-Coder to Understand Metastability or a Designer to Run Formal Methods

**CREATING AND EMBRACING GAPS:**
**INTENT VERIFICATION VS FUNCTIONAL VERIFICATION**
Design and DV were essentially the same skillset

- RTL, testbenches and tests were coded with same language
  - Directed, self-checking testing
  - Some smart checking of outputs
  - Manual maintenance of test coverage

- Engineered serially
  - Designed first, verified second
What drives verification today is different than design

**Design**
- Circuit Knowledge
- Sequential Function
- DFT
- RTL
- Assertions
- Object-Oriented

**Verification**
- Object-Oriented
- UVM
- UVVM/OSVVM
- Formal Methods
- Assertions
- Constrained-Random
- Functional Coverage
- Portable Stimulus
The Designer Knows the Intent: S1

But who is doing the verification?

- Scenario 1: Separate design and verification teams
  - Information is lost between design and verification
    - Subject to documentation, organization and knowledge gaps
  - Verification teams drive the tools but may not understand the goal
    - Should a verification engineer know about metastability?
**Scenario 2: One team of engineers**
- Engineers both design and verify
- Subject to skillset gaps
  - How many teams can afford engineers that are great at modern design *and* verification?
Why Not Let the Designer Verify?

Wrap verification tools in a designer-friendly wrapper

Intent Verification: Verification of design intent, by the engineer who knows the intent and correct implementation best

Create an intentional gap to embrace and align with your organizational & development gaps
Designer-Driven Intent Verification
Either way, designers need different tools

Static Creation
• Correct-by-construction
• Syntax

Static Lint
• Syntax
• Semantics
• Structural
• Standards

Formal Advanced Lint
• Sequential

Static & Formal CDC
• Clock Checks
• Protocol Verification

Static & Formal RDC
• Reset Checks
• Protocol Verification

HDL Designer Series

Questa® AutoCheck

Questa CDC

Questa RDC

Designers don’t want and shouldn’t have to use a testbench to verify intent
Why Isn’t Functional DV Enough?

• For functional DV to catch everything in “Intent Verification”, DV must:
  – Have full functional coverage, code coverage, testplan coverage, etc.
  – Sweep all asynchronous clocks and resets through all combinations
  – Model and detect failures correctly and completely

• Even if it does all that, it will take time to do so
  – Intent Verification solutions are static and exhaustive

• Even if it does all that, it could be debugged by another team
  – Even worse – it could be “fixed” by another team
Intent Verification Example: Lint

• Why wait for simulation to find an intent mismatch?

• The designer wants A=4
  – needs a check for missing parens

```verilog
wire [3:0] a, b, c, d;
assign b = 4’h2;
assign c = 4’h1;
assign a = 4’h8 >> b >> c;
```
Linting is Essential

We have a linting solution available

• Currently engaging customers in early access mode
• Implements important differentiators
• We are interested in working with additional customers
Intent Verification Example: Advanced Linting

- Why wait for simulation failures and debug (at best) to find:
  - A deadlock scenario in your state machine
  - An overflow condition on a registered variable
  - A combinational loop in your code, etc.

The designer can find and fix these things without a testbench, or knowledge of formal methods
Bridging the Gap to DV and Beyond

- Collateral from intent verification must be used in DV (and beyond)
  - Constraints/waivers should be valid and consistent through all Intent Verification
  - e.g. SDC for CDC & RDC, then on to synthesis, STA, etc.
Examples of Necessary Bridges to DV

When implementing an Intent Verification flow, consider the following:

• Some things are better left to simulation-based verification
  – Deep reconvergence verification for clock and reset domain checks
  – Some clock and reset domain crossing protocol verification

• Waiver cross-checking in simulation with functional coverage
  – Was the designer overly aggressive in waiving violations?

• But, verification environment must support variable latency and observation
Summary

• Identify gaps between design and verification

• Create an intentional gap by implementing Intent Verification flows
  – Embrace design/DV gap by aligning verification with organization/knowledge gaps

• Bridge the remaining gaps with cross-checks, DV and coverage
CLOSING ORGANIZATIONAL, GEOGRAPHICAL, TIME-BASED GAPS

SoCs Integrate IP from outside companies, across geographies, language barriers, previous projects
Recall the Story of a Very Large Aircraft...

A large aircraft launch was delayed multiple years

- Wire cables could not be connected in the prototype
  - Despite reviews, modeling and mock-ups

- Root-cause analysis determined the causes to be:
  - Different versions of software used by different teams in different countries
  - With different backgrounds, languages, goals, even measurement units

- **How different is that from your organization today?**

- **How can tools and methodologies bridge these gaps?**

  (Hint – minimize the handoffs)
Intent Verification, the SoC and Gaps

Flat Analysis
- Scalability (schedule) issues
- 3rd Party vendors don’t want to discuss low-level design details
- You should not need to re-verify IP that you’ve purchased or has already seen silicon (?)

Hierarchical Analysis
- Scalable (faster approach)
- Aligned with organizational gaps (IP vs. SoC, corporate, time)
- Requires abstraction
  - Black-box implies trust – nothing is re-evaluated beyond the lowest level
  - White-box will yield better integration analysis

Abstraction methods require EDA vendor support
Accuracy and performance both matter
Example: A Distributed Hierarchical CDC or RDC Approach

Enables a faster turnaround in case of RTL changes or ECOs
Hierarchical CDC or RDC Analysis Flow

**Top-down flow**
- Top-level setups
- Propagate top-level setups
- Top-level CDC
- Results

**Bottom-up flow**
- Block-level setups

**Steps**
1. Top-level CDC
2. Block-level CDC
3. Top-level CDC
4. Results
Conflicts will invalidate results!

- Critical for bottom-up flow
- Required for top-down flow with modified block constraints

1. Review block conflicts

2. Review top conflicts
Block-level Conflicts

- User vs Inferred port conflicts categorized in 3 sections
  1. Single-clock mismatch

  User: netlist port domain in -clock clk1 -module b1
  Inferred: hier port domain in -clock clk2 -module b1

  2. Multiple-clock mismatch

  User: netlist port domain in -clock clk1 -module b1
  Inferred: hier port domain in -multiple_clocks -module b1

  3. Sync mismatch

  User: netlist port domain in -async -module b1
  Inferred: hier port domain in -clock clk1 -module b1
Top-level Conflicts

- Check mismatch between Block and Top constraints
  - Constants & stable ports consistent
  - Clocks hooked up consistently
  - Ports hooked up consistently

- Warnings flagged for conflicts

<table>
<thead>
<tr>
<th>Port</th>
<th>Conflicting constraints</th>
</tr>
</thead>
</table>
| d1   | User-specified: netlist port domain d1 -clock clk1  
     |   Inferred: hier port domain d1 -clock clk2 |
| d2   | User-specified: hier port domain d2 -async  
     |   Inferred: netlist port domain d2[1] -async -clock clk1  
     |   hier port domain d2[0] -clock clk2 |
• 3rd Generation Hierarchical Data Model (HDM)
• Binary model
  – Database stores block information
  – Extendable to provide additional accuracy
  – Partially directives-based for overriding
• Supports both white-box & black-box abstractions
Questa HDM White-box Abstraction

• Crossings reported to/from internal HDM sequential elements
  – instead of HDM ports

• Improved multiple-fanout port visibility
  – Black-box: single crossing to/from the port reported for HDM
  – White-box: individual crossings for fanout registers of port’s load/driver

• Schematic visible for complete path including HDM internal logic

• Waivers work as expected
  – -through <hdm_port>
  – -to <hdm internal reg/latch>

• Enables sequential reconvergence verification
HDM Reconvergence : Example 1

- Syncs outside the HDM block converging inside the block

Reconvergence of synchronizers. (reconvergence)
------------------------------------------------------------------
clk2 : end : B1.q (/home/test01/dut.v : 1) (ID:reconvergence_79534)
clk2 : start : t2 (/home/test01/dut.v : 9) (Synchronizer ID:two_dff_19496) (Depth:0) (Reconvergence Severity:Caution)
clk2 : start : t4 (/home/test01/dut.v : 9) (Synchronizer ID:two_dff_67560) (Depth:0) (Reconvergence Severity:Caution)
HDM Reconvergence : Example 2

• Syncs inside the HDM block converging outside the block

Reconvergence of synchronizers. (reconvergence)
-----------------------------------------------------------------
clk2 : end : o1 (/home/test03/dut.v : 1) (ID:reconvergence_68716)
  clk2 : start : I1.q2 (/home/test03/dut.v : 13) (Depth:0) (Reconvergence Severity:Caution)
  clk2 : start : I1.r4 (/home/test03/dut.v : 15) (Depth:1) (Reconvergence Severity:Caution)
Hierarchical Analysis Results

- CDC and reconvergence checks run quickly
  - using HDMs on one of the SoCs
  - SoC contains 509 HDMs

These results clearly show the advantage of using HDMs.

<table>
<thead>
<tr>
<th>Turnaround Time / Memory</th>
<th>MPU (TOP)</th>
<th>IP1</th>
<th>IP2</th>
</tr>
</thead>
<tbody>
<tr>
<td>Register Count (M)</td>
<td>~.40M</td>
<td>~.16M</td>
<td>~.32M</td>
</tr>
<tr>
<td>Number HDM Blocks</td>
<td>509</td>
<td>2</td>
<td>1</td>
</tr>
<tr>
<td>CPU Run Time</td>
<td>~25Min</td>
<td>~17Min</td>
<td>~7 Hours</td>
</tr>
<tr>
<td>Max Memory</td>
<td>23.4 GB</td>
<td>4.3 GB</td>
<td>15.3 GB</td>
</tr>
</tbody>
</table>

- Reconvergence issue involves three IP blocks:
  - mod1, mod 2, and mod 3
  - HDMs encapsulate the schematic
  - Top-level shows the full reconvergence path
    - even when the three IP blocks are HDMs
It is probably not necessary to write an assertion that power is not ground.

Provavelmente não é necessário escrever uma afirmação de que o poder não está fundamentado.

It is probably not necessary to write a statement that power is not founded.
Can We Eliminate the Communication Gap?

Treat the Specification as the Golden Source

Specification-based tool flows fill the gap
e.g. for a CDC or RDC flow

• Specification is an important piece

• Guides both design and analysis

• Defines final closure requirements
Intent Verification: Specification Flow

• **Objective**
  – Specification to drive design and intent verification

• **Benefits**
  – Reduces the effort in results review and waiver specification

• **Use Model**
  – Specification constraints **guide** intent verification (e.g. CDC, RDC) analysis
  – Specification constraints **reduce** intent verification (e.g. CDC, RDC) violations
  – Designers do not review specification paths (correct by specification)

**Questa CDC supports a Specification Flow today**
Summary

• Projects are threatened by information loss across gaps

• Hierarchical abstraction bridges these gaps
  – Modeling must be accurate and provide high-visibility to be effective

• Documentation-rooted methodologies bridge these gaps
We Don’t Build RTL in Silicon. Verifying RTL is Not Enough.

CLOSING THE DESIGN/IMPLEMENTATION GAP
Intent Verification: Is Your Intent Preserved?

**RTL**

```vhdl
always @(posedge rx_clk)
begin: DMUX
    s1 <= tx_sel;
    rx_sel <= s1;
    if (rx_sel)
        rx_data <= tx_data1 && tx_data2;
    else
        rx_data <= rx_data ^ 4'b1111;
end
```

- Designer verifies intent at RTL (usually), but:
  - RTL mapped to logical equivalents in generic technology
  - Further optimized to target technology to meet constraints
- What is built in silicon is not what was verified in RTL!
  - This issue **should not** be ignored
You Verified CDC at RTL. ✔Intent, Right?

Did implementation introduce new CDC issues?
Synthesis Glitch Example: Functional

RTL Logic: Mux based synchronizer
Synthesis Glitch Example: Synthesized

For given constants, logic reduces to \((tx0 | \sim tx0)\) which causes glitch

Combo-logic implementation after synthesis
Synthesis Glitch Example: ECO

Adding the logic does not let the glitch propagate.
Manual Detection & Correction

• Detection:
  — Get lucky with gate-level simulations

• Correction:
  — Determine if asynchronous crossing, maps to RTL
    • name mappings, bit-blasting effects
  — Consider RTL constraint applications and waivers
  — Functionally analyze if glitch can occur
  — Add logic to the term to prevent the glitch
Closing the Design/Implementation Gap

How do you avoid falling in the gap?

• Prevention
  – Protect the DMUX paths
  – Guide the implementation

• Verification
  – Ensure signoff quality before tapeout
Prevention: Protecting DMUX Paths

• Infer the DMUX
  – Specify synthesis pragma to force DMUX implementation
    • Synthesis maps to a n-input mux or a tree of muxes

• Instantiate the DMUX from technology library
  – Inefficient for muxes with constant inputs
  – Need to have n-to-1 muxes and bus muxes

• Utilize synchronizer library
  – Protect the synch structures from synthesis optimization

Requires RTL changes and extra design review steps
Optimization can still move muxes (especially 2-to-1 mux) to LHS
Prevention: Guiding Synthesis

Use timing constraint
- to ensure the dmux is at the end of the fan-in cone to the RX register
- set_max_delay 0 -from <sel_sync_out> -to <rx_d>

Does not require RTL changes
Optimization can still move muxes (especially 2-to-1 mux) to LHS
These paths will be reported as timing violations
Verification: Questa Signoff CDC

**Setup**
- Gate Design
- Constraints at RTL

**CDC Analysis**
- Run CDC analysis
- Waivers at RTL
- CDC results
- Glitch results

**Debug**
- Fix CDC & glitch issues
- CDC sign-off

**Automated Setup**
- Constraints at RTL imported
- No iteration

**Reduced Noise**
- Bit-blasted paths regrouped
- Test logic automatically disabled

**Debug Eased**
- Waivers at RTL imported
- Correlation with RTL signals

**Glitches Detected**
- Structural & functional analysis of combinational logic for glitch
- Focused glitch debug
Glitch Check: Report

Glitch node and glitch signal
  • With clock group
  • Scheme name

Signal assignments exposing glitch

Paths with possible glitch
  • Could not prove with formal

Report logic under glitch condition
### Automotive Customer Results

- Two real glitches detected on the first design
- Verified 4 designs with Questa Signoff CDC
  - Design size ranging from ~1-10million

<table>
<thead>
<tr>
<th>Design</th>
<th>Gate count</th>
<th>Crossings</th>
<th>Glitches Detected</th>
<th>Formal (multicycle) proved glitches</th>
<th>Runtime (s)</th>
<th>Memory (GB)</th>
</tr>
</thead>
<tbody>
<tr>
<td>Design 1</td>
<td>~1M</td>
<td>3594</td>
<td>3</td>
<td>2</td>
<td>350</td>
<td>3.9</td>
</tr>
<tr>
<td>Design 2</td>
<td>~10M</td>
<td>297000</td>
<td>37</td>
<td>22</td>
<td>6661</td>
<td>40</td>
</tr>
<tr>
<td>Design 3</td>
<td>~5M</td>
<td>332258</td>
<td>313</td>
<td>20</td>
<td>2650</td>
<td>5.9</td>
</tr>
<tr>
<td>Design 4</td>
<td>~7M</td>
<td>128694</td>
<td>202</td>
<td>-</td>
<td>10310</td>
<td>23</td>
</tr>
</tbody>
</table>
Summary

• Gap: implementation process can introduce new issues
  – Intent verification flows on RTL are not enough

• Manual analysis flows require significant effort
  – Work through constraints, waivers, name mapping, bit-blasting, renaming
  – Must analyze if the issue merits an ECO

• Automated tool flows are better-suited
Verify that the product will work as intended

• Enable the designer to verify their intent was met, with an intent verification flow
• Align the verification tasks with organizational or skillset gaps

• Close the gap between what was verified and what is built

• Avoid lossy handoffs with strong and accurate modeling
• Automate generation and checking of constraints and assertions from documentation
We appreciate your attendance and hope you’ll watch your step

THANK YOU